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# AUTHENTICATION WITH AES CRYPTOGRAPHIC ENCRYPTION CHIP DESIGN AND SIMULATION

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#### ABSTRACT

Authentication is an indispensable part of Cryptography, which is an unconditionally secure key distribution technique based on the laws of nature. The paper explains the encryption technique using AES algorithm with 128 bits block size and 128 bit key size. The design is developed using S box technique with optimal hardware solution with the help of VHDL programming language. The design is synthesized on Vertex-5 FPGA using Xilinx 14.2 software and Modelsim 10.1 b. The algorithm is implemented to work in software and the file is to be encrypted in software and transferred to the machine containing the FPGA, which guarantees the secured communication and authentication. The concept of secured communication can be authenticated ATM machines, smart devices, biometrics applications and password protected devices. Remote authentication schemes allow a valid user to login a remote server and to access the services provided by the remote server over an insecure channel. Password based authentication schemes are the most widely used techniques for remote user authentication. Password based remote user authentication schemes are used to check the validity of a login request made by a remote user.

**KEYWORDS**: Advanced Encryption Standard (AES) method, Encryption and Decryption, Field programmable gate Array (FPGA).

### **INTRODUCTION**

Today's Cryptography is more applicable in secret writing, encryption and decryption techniques. Some of the more simple applications are secure communication, identification, authentication, and secret sharing. More complicated applications include systems for electronic commerce, certification, secures electronic mail, key recovery, secure computer access. When we send e-mail, we never know who know who would be reading or tampering it over the insecure channel? How can one keep message secret from all but those who are authorized to see it? How can the sender ensure himself that the message reaches the intended recipient unaltered by any unauthorized or unknown means? How can the receiver ensure himself that be message is coming from the legimate sender exactly as it was transmitted? This is all about why authenticity is required along with secrecy. When anyone logs into a host computer, ATM banking or any other such type of terminal, how does the host know who he/she is? How does the host know whether he/she is a valid user of the system? Authentication is the process of verifying that someone or something are who they claim to be before they are granted access to protected resources. Cryptography allows people to carry over the confidence found in the physical world to the electronic world, thus allowing people to do business electronically, without worries of deceit and deception. Everyday hundreds of people interact electronically, whether it is through E-mail, e-commerce (business conducted over the internet), ATM machines, or cellular phone. Cryptography makes secure web sites and electronic safe transmissions possible. The perpetual increase of information transmitted electronically has lead to an increased relies on cryptography. Authentication is a simple function where one party presents a set of credentials to a system. If the credentials match a given set on the system, the system returns a value that represents authorization; otherwise it does not. The purpose of authentication is to verify that the specific information presented represents a request to be authentic from a specified entity [1, 2].One of the simplest and most convenient security a mechanism is the use of a password authentication scheme. Examples of password authentication applications include remote login systems, ATM, PDA, and database management systems, etc. To access these resources, each user should have an identifier (ID) and a password (PW). The ID and PW are maintained by the remote system. Authentication can be based upon what he or she is a biometric like a fingerprint or voiceprint or any combination of these. Password- a combination of characters has

been the traditional and widely accepted tool of authentication .A valid user enters his/her password and host confirm it from the user every time he /she tries to log in[2, 5].

#### **AES ALGORITHM**

AES is based on the Rijndael cipher and developed by two Belgian cryptographers, Vincent Rijmen and Joan Daemen. They submitted a proposal to NIST during the AES selection process. Rijndael is a family of ciphers with different key and block sizes. The Advanced Encryption Standard (AES) [2, 4, 11] is standardized by the U.S. National Institute of Standards and Technology (NIST) in 2001 and is a specification for the encryption of electronic data established. For AES, NIST selected three members of the Rijndael family, each with a block size of 128 bits, but three different key lengths: 128, 192 and 256 bits. U.S government has been adopted AES and is now accepted worldwide. It is the advanced version of DES algorithm developed in 1977. The algorithm used in AES is based on symmetric key approach in which same key is shared by both encrypting and decrypting. In the US, AES was announced by the NIST as U.S. FIPS PUB 197 (FIPS 197) on November 26, 2001. This announcement followed a five-year standardization process in which fifteen competing designs were presented and evaluated, before the Rijndael cipher was selected as the most suitable. After the approval of Secretary of Commerce, AES became effective as a federal government standard on May 26, 2002. AES is included in the ISO/IEC 18033-3 standard. AES is available in many different encryption packages. AES is the first publicly accessible and open cipher approved by the National Security Agency (NSA) for top secret information when used in an NSA approved cryptographic module. The criteria defined by selecting AES fall into three areas Security, Implementation and cost of the algorithm. The main emphasis was the security of the algorithm to focus on resistance of cryptanalysis attacks, implementation cost should be less so it can be used for small devices like smart cards. The AES algorithm is a private key block cipher. It encrypts data of block size 128 bits. It uses three key sizes, 128 bits, 192 bits and 256 bits in three versions. AES uses three different types of round operations. Table I shows the number of rounds in three versions of AES. But, in each version final round key is 128 bits.

| Cipher Key size | No. of rounds | Round Key Size |
|-----------------|---------------|----------------|
| 128 bits        | 10            | 128 bits       |
| 192 bits        | 12            | 128 bits       |
| 256 bits        | 14            | 128 bits       |

 Table I: Round key size and number of rounds in three versions of AES [11, 14]
 \$\$\$

The initialization is done by adding first round key (128 bits) with 128 bits plain text. In subsequent steps, the following transformations are done: sub bytes, mix columns, shift rows and add round key. The last round is different from the previous rounds as there is no mix columns transformation. The internal 128 bits data in AES [3, 10] are represented in the form of 4x4 square matrix containing elements of size 8 bits and named as state elements. The decryption process involves of the inverse steps, decryption round contains of: Inverse S-BOX used for Byte Substitution, Inverse Shift Rows, Add Round Key and Inverse Mix Columns. The round keys will be generated using a unit called the key generation unit. This unit will be generating 176, 208 or 240 bytes of round keys depending on the size of the used key. The Add Round Key adds the round key word with each column of state matrix. It is similar to mix columns; the Add Round Key precedes one column at a time. The most important in this transformation, that it includes the cipher key. The state column will get XOR with key which is generated by key generator and create another state as shown in fig. 1

| <b>S</b> 0 | <b>S</b> 4 | <b>S</b> 8 | <b>S</b> 12 |     | KO | K4         | K8  | K12 |   | S0' | S4'         | S8'  | \$12'        |
|------------|------------|------------|-------------|-----|----|------------|-----|-----|---|-----|-------------|------|--------------|
| <b>S</b> 1 | <b>S</b> 5 | <b>S</b> 9 | <b>S</b> 13 |     | K1 | K5         | K9  | K13 |   | S1' | S5'         | S9'  | \$13'        |
| <b>S</b> 2 | <b>S</b> 6 | S10        | <b>S</b> 14 | (+) | K2 | K6         | K10 | K14 | = | S2' | <b>S6</b> ' | S10' | <b>S</b> 14' |
| <b>S</b> 3 | <b>S</b> 7 | S11        | S15         |     | K3 | <b>K</b> 7 | K11 | K15 |   | S3' | <b>S</b> 7' | S11' | \$15'        |
|            | State      |            |             |     |    |            |     |     |   | l)  | Sta         | ate  |              |

Fig.1 AES encryption with shift row transformations [11, 16]

Replacement of DES, Excellent security, limited key size. It encrypts data of block size 128 bits. It uses three key sizes, 128 bits, 192 bits and 256 bits in three versions. AES uses three different types of round operations





Fig.3 Subbyte Transformation [11]

The first transformation is Subbytes transformation used for encryption and inverse Sub Bytes used for decryption. The transformation is a nonlinear byte substitution that operates independently on each byte of the State using a substitution table (S-box). In the method multiplicative inverse is taken in the finite field GF  $(2^8)$  and affine transform to do the Subbytes transformation. In the inverse affine transform values have to find for inverse Subbytes transformation then multiplicative inverse of that byte. In the fig 3. it is indicated, the usage of the transformation can be done. There are two hexadecimal digits x and y in one state element, the left digit (x) defines the row and the right digit (y) defines the column of the substitution table. The junction of these two digits is the new bytes. Inverse Subbytes transformation is a method of inverse of Subbytes transformation. It is also possible to find in the similar way only table which is used for mapping the byte is different. The Subbytes transformation is done through S-box technique. There are two techniques to perform substitutions, (i) using composite field arithmetic (ii) using S-BOX table.

In the AES algorithm, there are four steps like Sub Byte, Shift Row, Mix Column and Add Round Key in normal rounds. The design consideration highlights some following modifications:

- Exclusion of shift row transformation
- Pipelining for high throughput with pipelined structure [17]
- Design optimization to keep handy balance between throughput and Silicon area

#### Exclusion of Shift Row Transformation

It is just routing information and takes no place in the design. Exclusion of Shift Row is performed through calling required shifted element from the data matrix. It is used instead of calling element one by one sequentially orderly from the data matrix. Thus merging of the two steps SUB-BYTE and SHIFT ROW reduces one step. The Fig No. 3 shows how exclusion of Shift Row is performed. The 16 elements are stored sequentially after each round in a register file. In the design mux selection required shifted data elements can be called from the register file and put into the S-Box instead of calling one by one. Sub Byte transformation (S-box) can be implemented as a look-up table (LUT) as shown in Fig. 4. This is a more efficient method than directly implementing the multiplicative inverse operation followed by affine transformation.



Fig.4 (a) Row Transformation [15]



Fig.4 (b) Shift Method [14]



Fig.4 (c) LUT with matrix [14]

In the AES algorithm, the primary requirement is of high data rate. High data rate can be achieved with the help of pipelined architecture. Pipelining is not for whole rounds required for encryption/decryption but one main constraint is that pipelining is only possible within each round. Further round can start only when previous round is totally completed as input data of the next round solely depend on the output of the previous rounds. Therefore, the design mainly concentrates on pipelined architecture implementation and methodology of each round. The design is optimized to keep handy balance between throughput and silicon area. There are different types of hardware architectures for AES algorithms are possible. The best architecture is one which is having the best tradeoff between clock speed and silicon Area for higher throughput. S-Box [15] transformation uses least 16 cycles are required to process 16 elements through S Box, if it is assumed in each cycle one element is processed through. Each AES round is having 4 steps. Using pipelined structures within each round, total cycles required (for 10 round of AES Algorithm for 128 bit data & 128 bit key) is approx 220 cycles. Here silicon area is saved but losing high data throughput. In the extreme case, whole 10 rounds can be completed within 44 cycles using 16 SBox (LUT). In this case, though high throughputs are achieved, silicon area in FPGA by introducing 4 SBox (LUT). Here whole encryption / decryption round completes within approx 90 cycles.

#### **ENCRYPTOR ARCHITECTURE**

The overall pipelined architecture for AES Encryptor is shown in fig. 5. The architecture has main FSM controller associated with one register bank, multiplexer, S-Box, latched pipelined register, column module mixed, the key expander, and add round key module along with the same controller structure. The user has to enter the plain text which has to be encrypted and also the key by which he will be able to decrypt this encrypted text back to plain text. Once the user specifies the key and the input data, a key expander module starts expanding the inputted original key so that it is sufficient for all the rounds of encryption procedure.



Fig. 5 AES Encryptor Architecture

Encryption process starts, after the key expansion is over. First of All, in encryption process, a pre round has to be performed which XORs the original 128 bits of input data with the original 128 bits of key and the intermediate result is stored in the register bank. Once the pre round is over then the remaining rounds of operation are performed. These rounds processed in the following manner: Intermediate register bank is used to read the data applied to four 4x1 mux of 32 bits from where the selected data is feed to the input of S-box. The S box transforms the data to their corresponding transformed data and pass it to mix column stage for further processing. The Mix column stage gets these 32 bits of data and according to the algorithm multiplies the data with a standard matrix to produce a 32 bit output. Now this output is applied to the Add Round Key[16, 17] which also has the 32 bits input from the key memory. These two 32 bits inputs are XORed in this module and are getting stored in the intermediate Register Bank again for the next round. Therefore, with the help of pipeline structure, all the four column of the new intermediate matrix are obtained one after the other keeping all the modules of the design busy all the time. In the last round, Mix Column stage is skipped and the result from the Add Round Key is gets stored in the output cipher text Memory. The proper selection of the module and data path for a particular round is done by the FSM controller [3]. The Controller also controls the Key Scheduling module so that valid keys are called for the particular round. The whole procedure of encryption is done for 128 bits plain text and 128 bits key in 88 cycles excluding the 62 cycles for the key expansion module. But the key expansion overhead does not cause degrading of the performance because for a input of larger size, the same expanded keys are again and again.

#### METHODOLOGY

The methodology comprises the different steps which are carried out for the design and development of chip. The model is shown in fig.6



Fig. 6 Steps in methodology

- *Design Specification:* In design specifications the designer decides to develop the chip either in top down or bottom up approach. In bottom up approach the circuit is designed using micro module to form a design and top down approach a design is distributed in sub modules. In the multistage networks bottom up approach is utilized to implement the design.
- *Network Configuration:* The designer has to choose the cluster size of the design. In the case of particular network, the block size and key size may be dependent or variable.
- *HDL Modeling:* The designer has to understand the feasibility to design with the supporting languages such as VHDL, Verilog HDL, and System C etc. The designer also decides the modeling of chip and design constructs in data flow, behavioral and structural model.
- *Functional Simulation:* The designed modules are checked according to their functionality and test cases. The functional simulation depends on the test benches developed by the designer, clock frequency and reset circuitry.
- *Pre Synthesis:* Pre synthesis includes the RTL simulation, device synthesis report contains the summary of hardware parameters with combinational and sequential circuit. If the hardware utilization is greater than the 100 % for the configured device, the designer has to redesign and check for the optimized device and timing parameters.
- *Experimentation and FPGA Synthesis:* The experimental setup is arranged to check the functionality of chip with its compatibility and interfacing to FPGA board. The maximum support frequency of FPGA board is analyzed to check the data transfer rate. In the experimental setup inputs can be through switches, and output

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can be analyzed with the help of LEDs. There is inbuilt ADC and DAC in Virtex 5 FPGA to check the real time functionality.

- *Parameters Analysis:* The FPGA source and target parameters are analyzed with the help of optimized FPGA results. Static timing analysis and device utilization is also analyzed with minim speed grade and memory utilization.
- *Testing:* The synthesized results are tested for the different test cases and combination with the help of LUTs. In the multistage network the inlets and outlets functionality is checked for maximum combination, inert and interexchange communication. The developed chip is also tested for analog input given to FPGA and processed output signals on DSO. The signal characteristics are also tested with the FPGA device compatibility and display unit.
- *Verification:* The Design Under Test (DUT) is verified with timing parameters and test cases. Standard VHDL has all the features necessary to code randomization of stimulus and functional coverage, both are very important while verifying larger, system-level designs. Verification is used to describe testing of a group of logic using a test bench, implemented for every verification level.

#### **RESULTS AND DESCRIPTION**

The simulation results RTL view of the chip is a top view representation depicting its pins details and input/ output logic. The possible inputs and ouputs used in the development of the chip are represented with their RTL view.



Fig. 7 (a) RTL view of encryption (128 bits data)

In the fig. 7(b), text[127:0] is the input text and ctext[127:0] is the out of cipher text. Clock is the common input.

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Fig. 5 (b) Modelsim output waveform of encryption (128 bits)

The description of the pins in listed in table 2

| Pins  | Description   |
|-------|---|
| text  | Input text of the encryption end it can be of 'N' bits, simulation results show for 128 bits. It is of    |
|       | std_logic_vector type   |
| clk   | clk is used for synchronization and provide clock pulse to work digital logic at duty cycle               |
| reset | Reset is used to reset the logic circuitry and synchronized with clock pulse of rising edge               |
| ctext | Ciphex text is the text which is encrypted at the transmitting end. It can be any garbage value and it is |
|       | of std_logic_vector type  |

Device utilization report [3] gives the percentage utilization of device hardware for the chip implementation. Device hardware includes No of slices, No of flip flops, No of input LUTs, No. of bounded IOBs and No of gated clocks (GCLKs) used in the implementation of design. Timing details provides the information of delay, minimum period value, maximum frequency value, minimum input arrival time before clock and maximum output required time after clock. Total memory utilization value required to complete the design. The target device is: xc5vlx20t-2-ff323 synthesized with Virtex-5 FPGA. Table 3 lists the simulated values of the design.

| Device utilization         |                  |      | Timing details                           |             |
|----------------------------|------------------|------|--|-------------|
| Device Part                |                  |      | Timing Parameter                         | Utilization |
| Utilization                |                  |      |  |             |
| Number of slices           | 126 out of 12480 | 1%   | Minimum period                           | 1.119ns     |
| Number of slice flip flops | 184 out of 12480 | 1%   | Maximum frequency                        | 410.00 MHz  |
| Number of 4 input LUTs     | 130 out of 189   | 68%  | Minimum input arrival time before clock  | 1.892ns     |
| Number of bonded IOBs      | 138 out of 172   | 80 % | Maximum output required time after clock | 2.110ns     |
| Number of GCLKs            | 1 out of 32      | 3%   | Total memory usage                       | 155009 kB   |

Table 3 Utilization in AES Design

# CONCLUSIONS

Remote users can be authenticated using passwords, smart cards etc. The authentication is integrated with Cryptographic AES encryption and decryption logic. The AES algorithm for encryption logic is implemented by selection of key and block size bits (128 bits).For higher throughput, 16 S-Box is used completing whole processes. In future work can be carried out for encryption and decryption of video. The proposed implementation is more efficient and suitable for hardware critical applications, and authenticate secured transmission with password facility. An optimized coding for the implementation of Rijndael algorithm for 128 bits has been developed which results in a throughput of Virtex 5 FPGA device. Architectural and algorithmic optimization like pipelining, exclusion of shift row, on the fly round key generation has been incorporated in the design.

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